

UNIVERSAL ITERATIVE METHODS FOR COMPUTING GENERALIZED INVERSES

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ABSTRACT. In this paper we construct a few iterative processes for computing $\{1, 2\}$ inverses of a linear bounded operator, based on the hyper-power iterative method or the Neumann-type expansion. Under the suitable conditions these methods converge to the $\{1, 2, 3\}$ or $\{1, 2, 4\}$ inverses. Also, we specify conditions when the iterative processes converge to the Moore-Penrose inverse, the weighted Moore-Penrose inverse or to the group inverse. A few error estimates are derived. The advantages of the introduced methods over the Tanabe's method [16] for computing the reflexive generalized inverses are also investigated.

KEY WORDS: generalized inverses, Moore-Penrose inverse, hyper-power method, Neumann-type expansion

1. INTRODUCTION

Let X and Y be two finite dimensional complex Hilbert spaces and let $A : X \rightarrow Y$ be a linear operator. There are well-known properties of generalized inverses of A :

$$(1) \quad AXA = A, \quad (2) \quad XAX = X, \quad (3) \quad (AX)^* = AX, \quad (4) \quad (XA)^* = XA.$$

For a subset \mathcal{S} of the set $\{1, 2, 3, 4\}$, the set of operators obeying the conditions contained in \mathcal{S} is denoted by $A\{\mathcal{S}\}$. An operator in $A\{\mathcal{S}\}$ is called an \mathcal{S} -inverse of A and is denoted by $A^{(\mathcal{S})}$. In particular, for any A , the set $A\{1, 2, 3, 4\}$ consists of a single element, the Moore-Penrose inverse of A , denoted by A^\dagger [9]. The group inverse $A^\#$ is the unique operator which satisfies (1), (2) and

$$(5) \quad AA^\# = A^\#A.$$

Any element from the class $A\{1, 2\}$ is also called the reflexive generalized inverse of A .

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Let $A = PQ$ be a full-rank decomposition of A , where P and Q are two full rank linear operators. Then all $\{1, 2\}$ inverses of A can be represented in the form

$$(6) \quad X = W_1(QW_1)^{-1}(W_2P)^{-1}W_2 = W_1(W_2AW_1)^{-1}W_2,$$

where W_1 and W_2 are suitable chosen operators, such that QW_1 and W_2P are invertible [13].

The weighted Moore-Penrose inverse is investigated in [3], [12]. For the sake of completeness we restate here the main results of these papers. Let there be given positive-definite (and hermitian) operators M and N . For any operator A there exists the unique solution $X = A_{M,N}^\dagger \in A\{1, 2\}$ satisfying the following equations in X [3], [12]:

$$(3M) \quad (MAX)^* = MAX \quad (4N) \quad (XAN)^* = XAN.$$

If $A = PQ$ is a full rank factorization of A , then [12]:

$$(7) \quad A_{M,N}^\dagger = (QN)^*(Q(QN)^*)^{-1}((MP)^*P)^{-1}(MP)^*.$$

The hyper-power method of the order 2 dates back to the well-known paper of Schulz [14]. Altman devised the hyper-power method of any arbitrary order $q \geq 2$, for inverting a nonsingular bounded operator on a Banach space [1]. In [10] the convergence of the same method is proved under the condition which is weaker than the one assumed in [1], and some better error estimates are derived.

Zlobec [19] and Petryshyn [11] showed that the q th order hyper-power iterative method with $q \geq 2$ for the determination of inverses of nonsingular matrices and linear operators, can be generalized to the determination of the Moore-Penrose inverse of an arbitrary matrix, or a bounded linear operator with closed range.

Zlobec in [19] defined the following two hyper-power iterative methods of an arbitrary high order $q \geq 2$:

$$\begin{aligned} T_k &= I_X - Y_k A, \\ Y_{k+1} &= (I_X + T_k + \cdots + T_k^{q-1})Y_k, \quad k = 0, 1, \dots \\ T'_k &= I_Y - AY'_k, \\ Y'_{k+1} &= Y'_k(I_Y + T'_k + \cdots + T'^{q-1}_k), \quad k = 0, 1, \dots \end{aligned}$$

It is well-known [19], that if we take

$$Y_0 = Y'_0 = \alpha A^*, \quad 0 < \alpha \leq \frac{2}{\text{tr}(A^*A)},$$

then $\lim_{k \rightarrow \infty} Y_k = \lim_{k \rightarrow \infty} Y'_k = A^\dagger$.

The process which generates the sequence Y_k is more superior than the process which generates the sequence Y'_k in the case $m > n$ [6]. The hyper-power iterative method of the order 2 is studied in [15] in view of the singular value decomposition of a matrix. In [6] the hyper-power iterative method is adapted for computing $A^\dagger B$, where A and B are arbitrary complex matrices with equal number of rows.

The paper is organized as follows. In Section 2 we construct iterative methods for computing the reflexive generalized inverses of a linear operator. These methods are based on the hyper-power iterative methods. We select two arbitrary matrices and adequate initial values for these methods to generate different generalized inverses for the concerned operator. In Section 3 we give a few error estimates, look for the optimal value of the parameter α and show that the method is self-correcting. In Section 4 we develop analogous iterative methods which arise from the Neumann-type expansion and compare our method with the Tanabe's method [16]. Finally, we give several examples which illustrate our theory.

2. ITERATIVE METHODS

In the following lemma we introduce an improvement of the hyper-power iterative method, and construct iterative method which generates all of the reflexive generalized inverses.

Lemma 2.1. *Let $\text{rank}(A) = r \geq 2$ and W_1, W_2 are two arbitrary operators, such that W_2AW_1 is invertible operator. If $q \geq 2$ is an integer, then the following two iterative processes:*

$$Y_0 = Y'_0 = \alpha(W_2AW_1)^*, \quad 0 < \alpha \leq \frac{2}{\text{tr}((W_2AW_1)^*W_2AW_1)},$$

$$\begin{aligned} T_k &= I_X - Y_kW_2AW_1, \\ Y_{k+1} &= (I_X + T_k + \cdots + T_k^{q-1})Y_k, \\ X_{k+1} &= W_1Y_{k+1}W_2 \quad k = 0, 1, \dots \end{aligned}$$

$$\begin{aligned} T'_k &= I_Y - W_2AW_1Y'_k, \\ Y'_{k+1} &= Y'_k(I_Y + T'_k + \cdots + T'^{q-1}_k), \\ X'_{k+1} &= W_1Y'_{k+1}W_2 \quad k = 0, 1, \dots \end{aligned}$$

generate the class of $\{1, 2\}$ inverses of A .

Proof. Using the results from [19], we conclude

$$\lim_{k \rightarrow \infty} Y_k = \lim_{k \rightarrow \infty} Y'_k = (W_2AW_1)^\dagger = (W_2AW_1)^{-1}.$$

According to (6), it is obvious that

$$\lim_{k \rightarrow \infty} X_k = \lim_{k \rightarrow \infty} X'_k = X = W_1(W_2AW_1)^{-1}W_2 \in A\{1, 2\}. \quad \square$$

Therefore, we just formed two iterative processes for computing all of the $\{1, 2\}$ inverses of A . However, under the suitable conditions, we can get some iterative methods for computing $\{1, 2, 3\}$ or $\{1, 2, 4\}$ inverses, the Moore-Penrose inverse, the weighted Moore-Penrose inverse or the group inverse of A . For the sake of simplicity we use the following notation: $B = W_2AW_1$, $C = AW_1$, $D = W_2A$.

Theorem 2.1. *Let $\text{rank}(A) = r \geq 2$ and QW_1, W_2P be invertible operators.*

- (a) *If W_2 is an unitary operator with respect to the considered scalar product and*

$$0 < \alpha \leq \min \left\{ \frac{2}{\text{tr}(B^*B)}, \frac{2}{\text{tr}(C^*C)} \right\},$$

then $X_k \rightarrow X = W_1(AW_1)^\dagger \in A\{1, 2, 3\}$ as $k \rightarrow \infty$.

- (b) *If W_1 is an unitary operator with respect to the considered scalar product and*

$$0 < \alpha \leq \min \left\{ \frac{2}{\text{tr}(B^*B)}, \frac{2}{\text{tr}(D^*D)} \right\},$$

then $X'_k \rightarrow X = (W_2A)^\dagger W_2 \in A\{1, 2, 4\}$ as $k \rightarrow \infty$.

- (c) *If (a) and (b) are valid, then $X_k \rightarrow A^\dagger$.*
 (d) *If (b) is valid and $W_2 = P^*$, then $X'_k \rightarrow X = A^\dagger$.*
 (e) *If (a) is valid and $W_1 = Q^*$, then $X_k \rightarrow X = A^\dagger$.*
 (f) *If $W_1 = Q^*$ and $0 < \alpha \leq \frac{2}{\text{tr}(B^*B)}$,*
then $X_k \rightarrow X = Q^(W_2AQ^*)^{-1}W_2 \in A\{1, 2, 4\}$.*
 (g) *If $W_2 = P^*$ and $0 < \alpha \leq \frac{2}{\text{tr}(B^*B)}$,*
*then $X_k \rightarrow X = W_1(P^*AW_1)^{-1}P^* \in A\{1, 2, 3\}$.*
 (h) *If $W_1 = Q^*$, $W_2 = P^*$ and $0 < \alpha \leq \frac{2}{\text{tr}(B^*B)}$,*

then $X_k \rightarrow A^\dagger$.

- (i) *If $W_1 = (QN)^*$, $W_2 = (MP)^*$ and $0 < \alpha \leq \frac{2}{\text{tr}(B^*B)}$,*

then $X_k \rightarrow A^\dagger_{M,N}$.

- (j) *In the case $m = n$, $W_1 = P$, $W_2 = Q$ and $0 < \alpha \leq \frac{2}{\text{tr}(B^*B)}$,*

we get $X_k \rightarrow A^\#$.

Proof. (a) Obviously,

$$X_{k+1} = W_1Y_{k+1}W_2 = W_1(I_X + T_k + \cdots + T_k^{q-1})Y_kW_2.$$

Let $Z_k = Y_k W_2$. Then

$$Z_{k+1} = Y_{k+1} W_2 = (I_X + T_k + \cdots + T_k^{q-1}) Y_k W_2 = (I_X + T_k + \cdots + T_k^{q-1}) Z_k,$$

where $T_k = I_X - Y_k W_2 A W_1 = I_X - Z_k (A W_1)$. Since

$$Z_0 = Y_0 W_2 = \alpha W_1^* A^* W_2^* W_2 = \alpha W_1^* A^* = \alpha (A W_1)^*,$$

we have [19] $Z_k \rightarrow (A W_1)^\dagger$, as $k \rightarrow \infty$. This implies

$$T_k^l = (I_X - Z_k A W_1)^l \xrightarrow[k \rightarrow \infty]{} (I_X - (A W_1)^\dagger A W_1)^l, \text{ for } l = 1, 2, \dots$$

Since $((A W_1)^\dagger A W_1)^2 = (A W_1)^\dagger A W_1$, we get

$$\begin{aligned} T_k^l &\xrightarrow[k \rightarrow \infty]{} I_X - \binom{l}{1} (A W_1)^\dagger A W_1 + \binom{l}{2} (A W_1)^\dagger A W_1 + \cdots + (-1)^l \binom{l}{l} (A W_1)^\dagger A W_1 \\ &= I_X - (A W_1)^\dagger A W_1. \end{aligned}$$

Hence

$$\begin{aligned} Z_{k+1} &= (I_X + T_k + \cdots + T_k^{q-1}) Z_k \xrightarrow[k \rightarrow \infty]{} \\ &\left[I_X + (I_X - (A W_1)^\dagger A W_1) + \cdots + (I_X - (A W_1)^\dagger A W_1) \right] (A W_1)^\dagger = \\ &= (A W_1)^\dagger. \end{aligned}$$

Now, it follows that $X_k \rightarrow W_1 (A W_1)^\dagger = X$, as $k \rightarrow \infty$. Since $(A W_1 (A W_1)^\dagger)^* = A W_1 (A W_1)^\dagger$, we get $(A X)^* = A X$ and X is an $\{1, 2, 3\}$ inverse for A .

(b) Let W_1 be unitary. For the sequences

$$X'_{k+1} = W_1 Y'_{k+1} W_2 = W_1 Y'_k (I_Y + T'_k + \cdots + T_k'^{q-1}) W_2$$

and $Z'_k = W_1 Y'_k$ we have

$$Z'_{k+1} = W_1 Y'_{k+1} = W_1 Y'_k (I_Y + T'_k + \cdots + T_k'^{q-1}).$$

Since $T'_k = I_Y - W_2 A Z'_k$ and $Z'_0 = W_1 Y'_0 = \alpha W_1 W_1^* A^* W_2^* = \alpha (W_2 A)^*$, we use the method from (a) to conclude that $Z'_k \rightarrow (W_2 A)^\dagger$ and $X'_{k+1} \rightarrow (W_2 A)^\dagger W_2 = X$. Since $((W_2 A)^\dagger W_2 A)^* = (W_2 A)^\dagger W_2 A$, we get $(X A)^* = X A$, and consequently, X is $\{1, 2, 4\}$ inverse of A .

(c) Follows from (a) and (b).

(d) In (b) we obtain $X_k \rightarrow (W_2A)^\dagger W_2 \in A\{1, 2, 4\}$. We mention one important property of the Moore–Penrose inverse [2], [4]:

$$(8) \quad (UV)^\dagger = V^\dagger U^\dagger \iff U^\dagger UVV^*U^* = VV^*U^* \text{ and } VV^\dagger U^*UV = U^*UV.$$

We take $U = W_2P$ and $V = Q$ in the expression $AX = PQ(W_2PQ)^\dagger W_2$. Operator W_2P is invertible, and Q^\dagger is the right inverse of the full rank operator Q . So the right side of (8) is valid in this case. Now, we get

$$AX = PQQ^\dagger(W_2P)^{-1}W_2 = P(W_2P)^{-1}W_2.$$

Also $(AX)^* = W_2^*(P^*W_2^*)^{-1}P^*$. Now, if $W_2 = P^*$, we get $(AX)^* = AX$. On the other hand, if (b) is valid, then $X = (W_2A)^\dagger W_2$ is an $\{1, 2, 4\}$ inverse of A , and we immediately conclude $X = A^\dagger$.

(e) If W_2 is an unitary operator, we have that $\lim_{k \rightarrow \infty} X_k = X = W_1(AW_1)^\dagger \in A\{1, 2, 3\}$. Using $U = P$ and $V = QW_1$, we conclude that (8) is valid. Consequently $XA = W_1(PQW_1)^\dagger PQ = W_1(QW_1)^{-1}P^\dagger PQ = W_1(QW_1)^{-1}Q$. Also $(XA)^* = Q^*(W_1^*Q^*)^{-1}W_1^*$. Obviously, if $W_1 = Q^*$, we get $X = A^\dagger$.

(f), (g) Follows from (6) and the well-known results [13]:

the general solution of the equations (1), (2), (4) is given by

$$X = Q^*(QQ^*)^{-1}(W_2P)^{-1}W_2 = Q^*(W_2AQ^*)^{-1}W_2;$$

the general solution of the equations (1), (2), (3) is given by

$$X = W_1(QW_1)^{-1}(P^*P)^{-1}P^* = W_1(P^*AW_1)^{-1}P^*.$$

(h) Follows from $A^\dagger = Q^*(QQ^*)^{-1}(P^*P)^{-1}P^*$ [3], [13].

(i) Comparing (6) and (7), we conclude that $A_{M,N}^\dagger$ can be selected from the class $A\{1, 2\}$ using $W_1 = (QN)^*$ and $W_2 = (MP)^*$.

(j) Follows from $A^\# = P(QP)^{-2}Q$ [5]. \square

In the case $\text{rank}(A) = 1$ we can use the next known proposition [19]:

Proposition 2.1. *Let A be of the rank $r = 1$. Then*

$$A^\dagger = \frac{1}{\text{tr}(A^*A)}A^*.$$

Theorem 2.2. *Let $\text{rank}(A) = 1$ and QW_1, W_2P be invertible operators.*

(a) $X = W_1(AW_1)^\dagger \in A\{1, 2, 3\}$ is given by

$$X = \frac{1}{\text{tr}((AW_1)^*AW_1)}W_1(AW_1)^*.$$

(b) $Y = (W_2A)^\dagger W_2 \in A\{1, 2, 4\}$ is given by

$$Y = \frac{1}{\text{tr}((W_2A)^*W_2A)}(W_2A)^*W_2.$$

(c) $Z = W_1(W_2AW_1)^{-1}W_2 \in A\{1, 2\}$ is presented by

$$Z = \frac{1}{\text{tr}((W_2AW_1)^*W_2AW_1)}W_1(W_2AW_1)^*W_2.$$

Proof. (a) Since QW_1 is invertible and P is left invertible, we have that $AW_1 = PQW_1 \neq 0$, which implies $\text{rank}(AW_1) \geq 1$. Using $\text{rank}(AW_1) \leq \text{rank}(A) = 1$, we get $\text{rank}(AW_1) = 1$. Thus, from Proposition 2.1. we get

$$(AW_1)^\dagger = \frac{1}{\text{tr}((AW_1)^*AW_1)}(AW_1)^*.$$

Now we prove that $X = W_1(AW_1)^\dagger$ is an $\{1, 2, 3\}$ inverse of A . It is easy to see that $W_1(AW_1)^\dagger$ is an $\{2, 3\}$ inverse of A . In order to prove that $AXA = A$ holds, we use the property (8) with $U = P$ and $V = QW_1$. Thus

$$AXA = PQW_1(PQW_1)^\dagger PQ = PQW_1(QW_1)^{-1}P^\dagger PQ = PQ = A.$$

(b) In a similar way can be proved that $\text{rank}(W_2A) = 1$. We just need to prove that $AYA = A$. We also use (8) with $U = W_2P$ and $V = Q$. Then

$$AYA = PQ(W_2PQ)^\dagger W_2PQ = PQQ^\dagger(W_2P)^{-1}W_2PQ = PQ = A. \quad \square$$

3. ERROR BOUNDS

Since $X = W_1(W_2AW_1)^{-1}W_2$ and $X_k = W_1Y_kW_2$, we have

$$\|X - X_k\| \leq \|W_1\| \cdot \|(W_2AW_1)^{-1} - Y_k\| \cdot \|W_2\|.$$

So we just have to make bounds for $\|(W_2AW_1)^{-1} - Y_k\|$. Operator $B = W_2AW_1$ is invertible, so B^*B is invertible and positive-definite. The spectrum of B^*B is $\sigma(B^*B) = \{\lambda_1, \lambda_2, \dots, \lambda_r\}$ and we can take that

$$0 < \lambda_1 < \lambda_2 < \dots < \lambda_r.$$

We need some matrix norm $\|\cdot\|$ such that $\|T_0\| < 1$. Let $\|\cdot\|_{sp}$ be the spectral norm, i.e.

$$\|C\|_{sp} = \sqrt{\max \lambda(C^*C)},$$

where $\max \lambda(C^*C)$ denotes the greatest eigenvalue of C^*C .

Lemma 3.1. *Let $T_0 = I - \alpha B^* B$ and*

$$0 < \alpha \leq \frac{2}{\operatorname{tr}(B^* B)}.$$

If $\operatorname{rank}(A) \geq 2$ then $\|T_0\|_{sp} < 1$.

Proof. Since $T_0^* T_0 = I - 2\alpha B^* B + \alpha^2 (B^* B)^2$, we get

$$\lambda \in \sigma(B^* B) \iff 1 - 2\alpha\lambda + \alpha^2\lambda^2 = (\alpha\lambda - 1)^2 \in \sigma(T_0^* T_0).$$

So $(\alpha\lambda - 1)^2 < 1$ if and only if $\alpha\lambda < 2$ for all $\lambda \in \sigma(B^* B)$. We see that $\|T_0\|_{sp} < 1$ is valid for $0 < \alpha < \frac{2}{\lambda_r}$. Since $\operatorname{tr}(B^* B) = \sum_{j=1}^r \lambda_j > \lambda_r$, we get

$$\frac{2}{\operatorname{tr}(B^* B)} < \frac{2}{\lambda_r}. \quad \square$$

We use Lemma 3.1 and the error estimates for the norm $\|(W_2 A W_1)^{-1} - Y_k\|$, given in [10], [8], [7], to verify the next theorem:

Theorem 3.1. *Providing that assumptions of Theorem 2.1. are valid, we get:*

- (a) $\|X - X_k\|_{sp} \leq \frac{\|Y_k T_k\|_{sp}}{1 - \|T_k\|_{sp}} \|W_1\|_{sp} \|W_2\|_{sp},$
- (b) $\|X - X_k\|_{sp} \leq \|T_{k-1}\|_{sp}^{q-1} \frac{\|Y_{k-1} T_{k-1}\|_{sp}}{1 - \|T_{k-1}\|_{sp}} \|W_1\|_{sp} \|W_2\|_{sp},$
- (c) $\|X - X_k\|_{sp} \leq \|T_0\|_{sp}^{q^k} \frac{\|Y_0\|_{sp}}{1 - \|T_0\|_{sp}} \|W_1\|_{sp} \|W_2\|_{sp},$
- (d) $\|X - X_k\|_{sp} \leq \frac{\|Y'_k\|_{sp} \|T'_k\|_{sp}}{1 - \|T'_k\|_{sp}} \|W_1\|_{sp} \|W_2\|_{sp},$
- (e) $\|X - X_k\|_{sp} \leq \frac{\|Y'_k\|_{sp} \|T'_{k-1}\|_{sp}^q}{1 - \|T'_{k-1}\|_{sp}^q} \|W_1\|_{sp} \|W_2\|_{sp},$
- (f) $\|X - X_k\|_{sp} \leq \frac{\|Y'_k\|_{sp} \|T'_{k-1}\|_{sp} \|(T'_{k-1})^{q-1}\|_{sp}}{1 - \|T'_{k-1}\|_{sp}} \|W_1\|_{sp} \|W_2\|_{sp},$
- (f) *The order of convergence of the defined processes is q , i.e.*

$$\|X - X_{k+1}\|_{sp} = O(\|X - X_k\|_{sp}^q), \quad k \rightarrow \infty.$$

If W_1 and W_2 are unitary operators, then $X = A^\dagger$, by Theorem 2.1. So we immediately get the error bounds for $\|A^\dagger - X_k\|_{sp}$, since $\|W_1\|_{sp} = \|W_2\|_{sp} = 1$.

Now we look for the optimal value of α .

Proposition 3.1. *Let $\text{rank}(A) \geq 2$. In each case of Theorem 2.1., the optimal value of α is the upper bound for the interval given in that case.*

Proof. Since I and αBB^* are commuting selfadjoint operators, the proof is identical to the corresponding in [19, Proposition 1]. \square

Remark 3.1. It is well-known that the hyper-power method for computing the inverse of an invertible operator is self-correcting, but it is not self-correcting for computing generalized inverses [17], [18]. There is the well-known Zielke's iterative refinement process which solves the self-correcting problem. Namely, the iterative refinement for computing the Moore-Penrose inverse of A has the form:

$$\begin{aligned}\tilde{X}_k &= A^* X_k^* X_k X_k^* A^* \\ \tilde{T}_k &= I - \tilde{X}_k A \\ X_{k+1} &= (I + \tilde{T}_k + \cdots + \tilde{T}_k^{q-1}) \tilde{X}_k.\end{aligned}$$

This modification is not necessary in each step.

Our iterative method for computing the sequence Y_k is self-correcting, so this method is self-correcting for computing the Z_k and X_k . We do not need any iterative refinements. So we solve, on an elementary way, the self-correcting problem for the iterative computation of generalized inverses.

4. USING THE NEUMANN-TYPE EXPANSION

It is well-known [16], [19] that the q -th order hyper-power method generates the partial sums of the infinite series

$$\sum_{i=0}^{\infty} \left[(I - X_0 A)^i X_0 \right] \quad \text{or} \quad \sum_{i=0}^{\infty} \left[X_0 (I - A X_0)^i \right],$$

i.e.

$$X_k = \sum_{i=0}^{q^k-1} \left[(I - X_0 A)^i X_0 \right] \quad \text{or} \quad X_k = \sum_{i=0}^{q^k-1} \left[X_0 (I - A X_0)^i \right].$$

In the case $\rho(I - X_0 A) < 1$ the inverse A^{-1} of a nonsingular matrix admits the Neumann-type expansion [16]

$$A^{-1} = \sum_{i=0}^{\infty} \left[(I - X_0 A)^i X_0 \right].$$

Similarly, in the case $\rho(I - A X_0) < 1$

$$A^{-1} = \sum_{i=0}^{\infty} \left[X_0 (I - A X_0)^i \right].$$

Zlobec [19] shown that A^\dagger can be computed by means of the infinite series under the assumption $0 < \alpha \leq \frac{2}{\text{tr}(A^*A)}$.

Our strategy is to adapt the infinite series in order to compute $(W_2AW_1)^{-1}$. In this way, we develop corresponding iterative method for computing the reflexive generalized inverses $W_1(W_2AW_1)^{-1}W_2$. We determine conditions for the defined method to generate the $\{1, 2, 3\}$, $\{1, 2, 4\}$ inverses, the Moore-Penrose or the group inverse.

If $q \geq 2$, $\text{rank}(A) = r \geq 2$, $B = W_2AW_1$, $C = AW_1$, $D = W_2A$, the iterative processes, based on the Neumann-type expansion, are defined as follows:

$$\begin{aligned} X_0 &= X'_0 = \alpha B^*, & 0 < \alpha &\leq \frac{2}{\text{tr}(B^*B)}, \\ X_k &= W_1 \cdot \sum_{i=0}^{q^k-1} \left[(I - X_0B)^i X_0 \right] \cdot W_2, & \text{or} \\ X'_k &= W_1 \cdot \sum_{i=0}^{q^k-1} \left[X_0 (I - BX_0)^i \right] \cdot W_2, & k = 0, 1, \dots \end{aligned}$$

The following results, analogous to the results from Theorem 2.1. can be easily verified.

Theorem 4.1. *Let $\text{rank}(A) = r \geq 2$ and QW_1, W_2P be invertible operators.*

(a) *If W_2 is an unitary operator and*

$$0 < \alpha \leq \min \left\{ \frac{2}{\text{tr}(B^*B)}, \frac{2}{\text{tr}(C^*C)} \right\},$$

then $X_k \rightarrow X = W_1(AW_1)^\dagger$ as $k \rightarrow \infty$ and X is an $\{1, 2, 3\}$ inverse of A .

(b) *If W_1 is an unitary operator and*

$$0 < \alpha \leq \min \left\{ \frac{2}{\text{tr}(B^*B)}, \frac{2}{\text{tr}(D^*D)} \right\},$$

then $X'_k \rightarrow X = (W_2A)^\dagger W_2 \in A\{1, 2, 4\}$ as $k \rightarrow \infty$.

(c) *If (a) and (b) are valid, then $X_k \rightarrow A^\dagger$.*

(d) *If (b) is valid and $W_2 = P^*$, then $X'_k \rightarrow X = A^\dagger$.*

(e) *If (a) is valid and $W_1 = Q^*$, then $X_k \rightarrow X = A^\dagger$.*

(f) *If $W_1 = Q^*$ and $0 < \alpha \leq \frac{2}{\text{tr}(B^*B)}$,*

then $X_k \rightarrow X = Q^(W_2AQ^*)^{-1}W_2 \in A\{1, 2, 4\}$.*

- (g) If $W_2 = P^*$ and $0 < \alpha \leq \frac{2}{\text{tr}(B^*B)}$,
 then $X_k \rightarrow X = W_1(P^*AW_1)^{-1}P^* \in A\{1, 2, 3\}$.
- (h) If $W_1 = Q^*$, $W_2 = P^*$ and $0 < \alpha \leq \frac{2}{\text{tr}(B^*B)}$,
 then $X_k \rightarrow A^\dagger$.
- (i) If $W_1 = (QN)^*$, $W_2 = (MP)^*$ and $0 < \alpha \leq \frac{2}{\text{tr}(B^*B)}$,
 then $X_k \rightarrow A_{M,N}^\dagger$.
- (j) In the case $m = n$, $W_1 = P$, $W_2 = Q$ and $0 < \alpha \leq \frac{2}{\text{tr}(B^*B)}$,
 we get $X_k \rightarrow A^\#$.

Proof. Follows from

$$\begin{aligned}
 X_k &= W_1 \cdot \sum_{i=0}^{q^k-1} \left[(I - X_0B)^i X_0 \right] \cdot W_2 = \\
 &= W_1 \cdot \sum_{i=0}^{q^k-1} \left[(I - X_0W_2AW_1)^i X_0W_2 \right], \quad k = 0, 1, \dots \\
 \\
 X'_k &= W_1 \cdot \sum_{i=0}^{q^k-1} \left[X_0 (I - BX_0)^i X_0 \right] \cdot W_2 = \\
 &= \sum_{i=0}^{q^k-1} \left[W_1X_0 (I - W_2AW_1X_0)^i \right] \cdot W_2, \quad k = 0, 1, \dots \quad \square
 \end{aligned}$$

Remark 4.1. Now we give some comparisons with the paper of Tanabe [16]. In [16] there are given necessary and sufficient conditions for the starting approximation X_0 of the hyper-power iterative method or the Neumann-type series, ensuring convergence of these methods to an arbitrary reflexive generalized inverse. Advantages of this paper related to [16] are:

- (a) It is more convenient to use the initial conditions in Theorem 2.1. than the conditions from [16, Theorem 2.1].
- (b) We give a few error estimates.
- (c) We know exact conditions ensuring convergence of defined processes to the $\{1, 2, 3\}$ inverses, the $\{1, 2, 4\}$ inverses, the Moore-Penrose inverse, the group inverse or the weighted Moore-Penrose inverse of A .

5. EXAMPLES

Example 5.1. Consider the matrix $A = \begin{pmatrix} 1 & 0 \\ -1 & 0 \\ 0 & 1 \end{pmatrix}$. If we select unitary matrix $W_1 = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$, and $W_2 = \begin{pmatrix} 1 & 0 & 1 \\ 0 & 1 & 0 \end{pmatrix}$, we get

$$B = W_2 A W_1 = \begin{pmatrix} 1 & 1 \\ 0 & -1 \end{pmatrix}, \quad \text{tr}(B B^*) = 3,$$

$$D = W_2 A = \begin{pmatrix} 1 & 1 \\ -1 & 0 \end{pmatrix}, \quad \text{tr}(D D^*) = 3,$$

$$\alpha = \min \left\{ \frac{2}{3}, \frac{2}{3} \right\}.$$

Using the package MATHEMATICA we construct the following iterative process of the order 4:

$$Y_0 = \alpha B^* = \begin{pmatrix} \frac{2}{3} & 0 \\ \frac{2}{3} & -\frac{2}{3} \end{pmatrix},$$

$$T_0 = I - B Y_0 = \begin{pmatrix} -\frac{1}{3} & \frac{2}{3} \\ \frac{2}{3} & \frac{1}{3} \end{pmatrix},$$

$$Y_1 = Y_0(I + T_0 + T_0^2 + T_0^3) = \begin{pmatrix} \frac{56}{81} & \frac{56}{81} \\ 0 & -\frac{56}{81} \end{pmatrix}.$$

$$T_1 = I - B Y_1 = \begin{pmatrix} \frac{25}{81} & 0 \\ 0 & \frac{25}{81} \end{pmatrix},$$

$$Y_2 = Y_1(I + T_1 + T_1^2 + T_1^3) = \begin{pmatrix} \frac{42656096}{43046721} & \frac{42656096}{43046721} \\ 0 & -\frac{42656096}{43046721} \end{pmatrix}.$$

In a similar way can be obtained

$$Y_3 = \begin{pmatrix} \frac{3433683797009448119270886198656}{3433683820292512484657849089821} & \frac{3433683797009448119270886198656}{3433683820292512484657849089821} \\ 0 & -\frac{3433683797009448119270886198656}{3433683820292512484657849089821} \end{pmatrix}.$$

Now we get the following sequence $X_k = W_1 Y_k W_2$:

$$X_1 = \begin{pmatrix} 0 & -\frac{56}{81} & 0 \\ \frac{56}{81} & \frac{56}{81} & \frac{56}{81} \end{pmatrix},$$

$$X_2 = \begin{pmatrix} 0 & -\frac{42656096}{43046721} & 0 \\ \frac{42656096}{43046721} & \frac{42656096}{43046721} & \frac{42656096}{43046721} \end{pmatrix},$$

$$X_3 = \begin{pmatrix} 0 & -\frac{3433683797009448119270886198656}{3433683820292512484657849089821} & 0 \\ \frac{3433683797009448119270886198656}{3433683820292512484657849089821} & \frac{3433683797009448119270886198656}{3433683820292512484657849089821} & \frac{3433683797009448119270886198656}{3433683820292512484657849089821} \end{pmatrix}.$$

We have obtained a sequence converging to $X = \begin{pmatrix} 0 & -1 & 0 \\ 1 & 1 & 1 \end{pmatrix} \in A\{1, 2, 4\}$.

Example 5.2. For the same matrices A , W_1 and $W_2 = P^* = \begin{pmatrix} 1 & -1 & 0 \\ 0 & 0 & 1 \end{pmatrix}$, we get

$$\begin{aligned} B &= W_2 A W_1 = \begin{pmatrix} 0 & 2 \\ 1 & 0 \end{pmatrix}, & \operatorname{tr}(B B^*) &= 5, \\ D &= W_2 A = \begin{pmatrix} 2 & 0 \\ 0 & 1 \end{pmatrix}, & \operatorname{tr}(D D^*) &= 5, \\ \alpha &= \min \left\{ \frac{2}{5}, \frac{2}{5} \right\}. \end{aligned}$$

The following sequence $X_k = W_1 Y_k W_2$ can be obtained:

$$\begin{aligned} X_1 &= \begin{pmatrix} \frac{272}{625} & -\frac{272}{625} & 0 \\ 0 & 0 & \frac{544}{625} \end{pmatrix}, \\ X_2 &= \begin{pmatrix} \frac{76272421952}{152587890625} & -\frac{76272421952}{152587890625} & 0 \\ 0 & 0 & \frac{152544843904}{152587890625} \end{pmatrix}, \end{aligned}$$

The first row of X_3 is

$$\begin{array}{ccc} \frac{271050543121374391659953053961243098931900672}{542101086242752217003726400434970855712890625} & -\frac{271050543121374391659953053961243098931900672}{542101086242752217003726400434970855712890625} & 0 \end{array}$$

and the second row of X_3 is

$$0 \quad 0 \quad \frac{542101086242748783319906107922486197863801344}{542101086242752217003726400434970855712890625}.$$

The limit of this sequence is $X = \begin{pmatrix} \frac{1}{2} & -\frac{1}{2} & 0 \\ 0 & 0 & 1 \end{pmatrix} = A^\dagger$.

Example 5.3. For the same matrix A and $W_1 = Q^* = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}$, $W_2 = P^* = \begin{pmatrix} 1 & -1 & 0 \\ 0 & 0 & 1 \end{pmatrix}$, we can generate the same sequence X_k as in Example 5.2.

Example 5.4. Now we expand Example 5.1. to include illustrations for the values of α less than the optimal, or for values which do not produce the convergence.

Modifying only $\alpha = 1/2 < 2/3$ in Example 5.1., we obtain the following sequence:

$$\begin{aligned} X_1 &= \begin{pmatrix} \frac{3}{16} & -\frac{11}{16} & \frac{3}{16} \\ \frac{11}{16} & \frac{1}{2} & \frac{11}{16} \end{pmatrix}, \\ X_2 &= \begin{pmatrix} \frac{987}{65536} & -\frac{63939}{65536} & \frac{987}{65536} \\ \frac{63939}{65536} & \frac{7869}{8192} & \frac{63939}{65536} \end{pmatrix}, \\ X_3 &= \begin{pmatrix} \frac{10610209857723}{18446744073709551616} & -\frac{18446726906029374051}{18446744073709551616} & \frac{10610209857723}{18446744073709551616} \\ \frac{18446726906029374051}{18446744073709551616} & \frac{2305839536977439541}{2305843009213693952} & \frac{18446726906029374051}{18446744073709551616} \end{pmatrix}. \end{aligned}$$

The obtained sequence converges to the same matrix as in Example 5.1., but the convergence is slower.

Similarly, using $\alpha = 1/3 < 1/2 < 2/3$ we obtain the following:

$$X_1 = \begin{pmatrix} \frac{7}{27} & -\frac{47}{81} & \frac{7}{27} \\ \frac{47}{81} & \frac{26}{81} & \frac{47}{81} \end{pmatrix},$$

$$X_2 = \begin{pmatrix} \frac{726103}{14348907} & -\frac{39522143}{43046721} & \frac{726103}{14348907} \\ \frac{39522143}{43046721} & \frac{37343834}{43046721} & \frac{39522143}{43046721} \end{pmatrix},$$

$$X_3 = \begin{pmatrix} \frac{83909608561183162716808087}{1144561273430837494885949696427} & \frac{3433276514496608404104017015327}{3433683820292512484657849089281} & \frac{83909608561183162716808087}{1144561273430837494885949696427} \\ \frac{3433276514496608404104017015327}{3433683820292512484657849089281} & \frac{3433024785670924854615866591066}{3433683820292512484657849089281} & \frac{3433276514496608404104017015327}{3433683820292512484657849089281} \end{pmatrix}$$

The values of α , greater than the optimal one leads to the divergence of the methods. For example, $\alpha = 1 > 2/3$ imply the following divergent sequence:

$$X_1 = \begin{pmatrix} -3 & 1 & -3 \\ -1 & 2 & -1 \end{pmatrix},$$

$$X_2 = \begin{pmatrix} -987 & 609 & -987 \\ -609 & 378 & -609 \end{pmatrix},$$

$$X_3 = \begin{pmatrix} -10610209857723 & 6557470319841 & -10610209857723 \\ -6557470319841 & 4052739537882 & -6557470319841 \end{pmatrix}.$$

Finally, a small decreasing of the parameter α near the optimal value $\frac{2}{\text{tr}(B^*B)}$ imply slightly slowing down the speed of the convergence. For example, $\alpha = 2/3 - 0.001$ generates the following sequence:

$$X_1 = \begin{pmatrix} 0.00221225 & -0.692092 & 0.00221225 \\ 0.692092 & 0.68988 & 0.692092 \end{pmatrix},$$

$$X_2 = \begin{pmatrix} 0.000255561 & -0.991009 & 0.000255561 \\ 0.991009 & 0.990753 & 0.991009 \end{pmatrix},$$

$$X_3 = \begin{pmatrix} 7.1253910^{-10} & -1. & 7.1253910^{-10} \\ 1. & 1. & 1. \end{pmatrix}.$$

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